# Exact values for Waring's problem in finite fields

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## Waring's problem in finite fields

joint work with Arne Winterhof

Over a ring R, Waring's problem in degree n asks whether every element a of R can be written in the form

$$a = \sum_{i} a_{i}^{n} \tag{1}$$

for some  $a_i \in R$ , and whether the number of terms needed can be uniformly bounded for all  $a \in R$ .

The problem is best known over  $\mathbb{Z}$ , but was also much studied in the case where R is a finite field (see Winterhof (1998) for a survey).

We define the Waring function g(k,q) as follows: if all  $a \in \mathbb{F}_q$  have an expansion (1), then g(k,q) is the maximal number of terms needed for any a; otherwise, g(k,q) is undefined.

### Some results on the Waring function

- We may assume that the exponent k divides q-1.
- If  $k^2 < q$  or if q is prime, then g(k,q) exists.
- A counterexample: q nonprime, k = q 1.
- If g(k,q) exists, then  $g(k,q) \leq k$  (inhomogeneous Chevalley-Warning); there is then a deterministic polynomial time algorithm to solve

$$a_1^k + \ldots + a_k^k = a.$$

• If  $(k-1)^4 < q$ , then g(k,q) = 1 or 2 (Weil bound). Assuming this, in fact, whenever  $abc \neq 0$ , then

$$ax^k + by^k = c$$

is solvable.

### **Reduction to the prime field**

We have the following nice inequality: if  $g(k, p^n)$  exists, then

 $g(k, p^n) \le ng(d, p),$ 

with  $d = \frac{k}{\gcd(k, \frac{p^n-1}{p-1})}$ .

This follows because  $g(k, p^n)$  exists if and only if

$$\mathbb{F}_{p^n} = \mathbb{F}_p(\alpha^k)$$

for some  $\alpha \in \mathbb{F}_{p^n}$ , so we have

$$a = a_0 + a_1 \alpha^k + \ldots + a_{n-1} \alpha^{(n-1)k},$$

and we write each  $a_i$  as a sum of dth powers in  $\mathbb{F}_p$ . Finally, more elements of  $\mathbb{F}_p$  may become dth powers in the extension field.

We use this reduction in the sequel.

#### **Basic setup and results**

We have odd primes p and r, with p a primitive root modulo r. Thus,

 $\mathbb{F}_{p^{r-1}} \text{ is generated over } \mathbb{F}_p \text{ by } \zeta_r.$ We let  $k = \frac{p^{r-1}}{r}$  or  $\frac{p^{r-1}}{2r}$ , so *k*th powers are *r*th or 2*r*th roots of unity. We compute  $g(k, p^{r-1})$  for these cases:

Theorem We have 
$$\begin{cases} g\left(\frac{p^{r-1}-1}{r}, p^{r-1}\right) = \frac{(p-1)(r-1)}{2}.\\\\ g\left(\frac{p^{r-1}-1}{2r}, p^{r-1}\right) = \begin{cases} \left\lfloor \frac{pr}{4} - \frac{p}{4r} \right\rfloor & \text{if } r < p;\\\\ \left\lfloor \frac{pr}{4} - \frac{r}{4p} \right\rfloor & \text{if } r \ge p. \end{cases}\end{cases}$$

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# Norm and weight

Let's call the minimal number of k'th powers needed to form a the Waring weight of a. For a vector  $\mathbf{a} \in \mathbb{F}_p^r$ , we write

$$\|\mathbf{a}\| = |a_0| + \ldots + |a_{r-1}|,$$

where |a| is the Waring weight of a (depending on the exponent k).

Lemma Let  $a = a_0 + a_1\zeta_r + \ldots + a_{r-1}\zeta_r^{r-1}$ . Then the weight of a is equal to

$$\min\{\|\mathbf{a} + x\mathbf{e}\| : x \in \{0, 1, \dots, p-1\}\}.$$

Lemma If  $k = \frac{p^{r-1} - 1}{r}$ , then  $|a| = a^r$  for all  $a \in \mathbb{F}_p$ .

Lemma If  $k = \frac{p^{r-1}-1}{2r}$ , then  $|a| = \text{``min}\{a, p-a\}$ '' for all  $a \in \mathbb{F}_p$ .

# **Reformulation of the problem**

We can now reformulate the Waring problem for these cases as follows.

Let  $V = (\mathbb{Z}/m\mathbb{Z})^r$ , let  $|\cdot|: \mathbb{Z}/m\mathbb{Z} \to \mathbb{Z}$  be some weight function on  $\mathbb{Z}/m\mathbb{Z}$ , and for  $\mathbf{a} \in V$ , let  $||\mathbf{a}|| = |a_0| + \ldots + |a_{r-1}|$ .

We say  $\boldsymbol{a}$  is admissible if

 $\|\mathbf{a}\| \leq \|\mathbf{a} + x\mathbf{e}\|$  for all  $x \in \mathbb{Z}/m\mathbb{Z}$ .

Now we want to know the maximal norm of an admissible vector.

We use the weights defined earlier, i.e.,

 $|a|_1$  is the smallest nonnegative integer representative of a;

 $|a|_2$  is the absolute value of the symmetric integer representative of a.

# An upper bound on $\|\mathbf{a}\|$

If  $\|\mathbf{a}\|_i \leq \|\mathbf{a} + x\mathbf{e}\|_i$  for all x, add these and get

$$m \|\mathbf{a}\|_{i} \le r \sum_{x=0}^{m-1} |x|_{i}.$$
 (2)

Lemma We have

$$\sum_{x} |x|_{1} = \frac{m(m-1)}{2} \quad ; \quad \sum_{x} |x|_{2} = \begin{cases} \frac{m^{2}}{4} & \text{if } m \text{ is even} \\ \frac{m^{2}-1}{4} & \text{if } m \text{ is odd.} \end{cases}$$

We refine (2) a little bit by noting that

$$\|\mathbf{a} + x\mathbf{e}\|_{1} \equiv \|\mathbf{a}\| + rx \pmod{m},$$
  
so in fact we have  $\|\mathbf{a}\|_{1} \leq \|\mathbf{a} + x\mathbf{e}\|_{1} - |rx|_{1}$ , and get the sharp  
 $\|\mathbf{a}\|_{1} \leq \frac{mr - m - r + \gcd(m, r)}{2}.$ 

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#### An aside, with an open problem

In general, assume q is a positive integer and  $p \equiv 1 \pmod{q}$  is prime. Let  $\zeta_q$  be a primitive qth root of unity in  $\mathbb{F}_p$ , and define

$$|x|_q = \min\{ |\zeta^i x|_1 : 0 \le i \le q - 1 \}.$$

Note that this agrees with our earlier definition of  $|\cdot|_2$ .

Proposition We have for  $q \ge 2$ 

$$\sum_{x} |x|_q = \left(\frac{1}{q+1} - \frac{B_q}{q!}\right) p^2 + O(p^{2-\varepsilon}),$$

where  $B_q$  is the *q*th Bernoulli number.

Conjecture We have

$$\sum_{x} |x|_{3} = \frac{p^{2} - 1}{4}.$$

Any takers??

## An upper bound for ||a|| (continued)

Recall that  $\mathbf{a} \in V = (\mathbb{Z}/m\mathbb{Z})^r$ , and  $|x|_2 = \min\{x, m - x\}$ .

For  $|\cdot|_2$ , the upper bound on  $||\mathbf{a}||_2$  for admissible vectors  $\mathbf{a}$  that we get is sharp whenever  $r \ge m$  or r is even. If r < m and r is odd, we consider the norm sequence

$$N_x = \|\mathbf{a} + x\mathbf{e}\|,$$

and using symmetry properties of this sequence, we derive a sharp bound in this case also.

We have, for admissible  $\mathbf{a} \in V$ ,

$$\|\mathbf{a}\|_{2} \leq \begin{cases} \frac{mr}{4} & \text{if } m \text{ and } r \text{ are even;} \\ \left\lfloor \frac{mr}{4} - \frac{1}{2} \right\rfloor & \text{if } m \text{ is even, } r \text{ is odd, and } r > m; \\ \left\lfloor \frac{mr}{4} - \frac{r}{4m} \right\rfloor & \text{if } m \text{ is odd and } r \ge m; \\ \left\lfloor \frac{mr}{4} - \frac{1}{2} \right\rfloor & \text{if } m \text{ is odd, } r \text{ is even, and } r < m; \\ \left\lfloor \frac{mr}{4} - \frac{m}{4r} \right\rfloor & \text{if } r \text{ is odd and } r < m. \end{cases}$$

# Matching up

To show that the given upper bounds are sharp, we need to construct admissible vectors attaining the bound.

If m and r are even,  $(0, \ldots, 0, \frac{m}{2}, \ldots, \frac{m}{2})$  is admissible of norm mr/4, which is maximal.

If m is odd and r is even, we use  $(0, \frac{m-1}{2})$  as a building block, with some cunning.

For odd r, the constructions are rather involved. First, by induction we reduce to the case that r < 2m. Then, we solve some integer programming problems with the goal to make the norm sequence, which has  $N_{x+1} \neq N_x$  for all x, as smooth and as flat as possible. Finally, the case of odd m is derived from the case of even m.

### Recapitulation

Theorem Let p and r be odd primes, with p a primitive root modulo r. Then we have

$$g\left(\frac{p^{r-1}-1}{r}, p^{r-1}\right) = \frac{(p-1)(r-1)}{2}.$$
$$g\left(\frac{p^{r-1}-1}{2r}, p^{r-1}\right) = \begin{cases} \left\lfloor \frac{pr}{4} - \frac{p}{4r} \right\rfloor & \text{if } r < p;\\ \left\lfloor \frac{pr}{4} - \frac{r}{4p} \right\rfloor & \text{if } r \ge p. \end{cases}$$

Furthermore, there exists an algorithm that shows elements in  $\mathbb{F}_{p^{r-1}}$  that need this many terms when writing them as sum of *k*th powers (KASH 2.5 code available...).

Note that all bounds are symmetric in p and r!