

**Operations Research**  
**winter term 2019/2020**

**1st work sheet (integer programming)**

1. As the leader of an oil-exploration drilling venture you must determine the least-cost selection of 5 out of 10 possible sites. Label the sites  $S_1, S_2, \dots, S_{10}$  and the corresponding exploration costs by  $C_i$ ,  $1 \leq i \leq 10$ , respectively. Consider the following regional development restrictions.

- (i) Evaluating sites  $S_1$  and  $S_7$  will prevent you from exploring site  $S_8$ .
- (ii) Evaluating site  $S_3$  or  $S_4$  prevents you from assessing site  $S_5$ .
- (iii) Of the group  $S_5, S_6, S_7$  and  $S_8$  only two sites may be assessed.

Formulate an integer program to determine the minimum cost exploration scheme that satisfies these restrictions.

2. Three different items are to be routed through three machines. Each item must be processed first on machine 1, then on machine 2 and finally on machine 3. The sequence of items may differ for each machine. Assume that for  $i, j \in \{1, 2, 3\}$  the times  $t_{ij}$  required to perform the work on item  $i$  by machine  $j$  are known and are integers. Our objective is to minimize the total time necessary to process all the items.

- (a) Formulate the problem as an integer programming problem.  
Hint: Let  $x_{ij}$  be the starting time of processing item  $i$  on machine  $j$ . Your model must prevent two items from occupying the same machine at the same time; also an item may not start processing on machine  $j + 1$  unless it has completed processing on machine  $j$ .
- (b) Suppose we want the items to be processed in the same sequence on each machine. Change the formulation in part (a) accordingly.

3. Consider a placement problem for firehouse locations. Assume that the population is concentrated in  $I$  districts within the city and that for each  $i \in I$  district  $i$  contains  $p_i$  people. Preliminary analysis has limited the potential location of firehouses to  $J$  sites. Let  $d_{ij}$  be the distance from the center of district  $i$  to site  $j$ , for  $i \in I, j \in J$ . The goal is to determine the “best” site selection and the assignment of the districts to firehouses such that the following criteria are fulfilled.

- (i) Every district is assigned to at least one firehouse.
- (ii) It costs  $f_j(s_j)$  to build a firehouse at site  $j$  to service  $s_j$  people; for each  $j \in J$  the function  $f_j$  is convex and piecewise linear with two linear pieces with gradients  $a_j$  and  $b_j$ ,  $a_j < b_j$ , respectively.
- (iii) There is a total budget  $B$  allocated for the firehouse construction.
- (iv) A central district is particularly susceptible to fire and either sites 1 and 2 or sites 3 and 4 can be used to protect this district.
- (v) The social-welfare goal is to minimize the distance travelled to the district farthest from its assigned firehouse.

Formulate this problem as a linear integer program.

4. In the lecture we discussed the following simplified version of the warehouse location problem.

$$\begin{aligned} \text{minimize} \quad & \sum_i \sum_j c_{ij} x_{ij} + \sum_i f_i y_i, \\ \text{subject to} \quad & \sum_i x_{ij} = d_j, \text{ for } j = 1, 2, \dots, n \end{aligned}$$

$$\sum_j x_{ij} - y_i \sum_j d_j \leq 0, \text{ for } i = 1, 2, \dots, m$$

$$x_{ij} \geq 0, \text{ for } i = 1, 2, \dots, m, j = 1, 2, \dots, n$$

$$y_i \in \{0, 1\}, \text{ for } i = 1, 2, \dots, m.$$

- (a) The above model assumes only one product and two distribution stages (from warehouse to customer). Indicate how the model would be modified if there were three distribution stages (from plant to warehouse to customer) and  $L$  different products.  
Hint: Define a new decision variable  $x_{ijkl}$  equal to the amount to be sent from plant  $i$ , through warehouse  $j$ , to customer  $k$ , of product  $l$ .
- (b) Suppose there are maximum capacities for plants and size limits, both upper and lower bounds, for warehouses. What is the relevant model now?
- (c) Assume that each customer has to be served from a single warehouse, i.e., no splitting of order is allowed. How should the warehousing model be modified?
- (d) Assume that each warehouse  $i$  experiences economies of scale when shipping above a certain threshold quantity to an individual customer  $j$ ; i.e., the unit distribution cost is  $c_{ij}$  whenever the amount shipped is between 0 and  $d_{ij}$ , and  $c'_{ij}$  (lower than  $c_{ij}$ ) whenever the amount shipped exceeds  $d_{ij}$ . Formulate the warehouse location model under these conditions.
5. Consider the following word game. You are assigned a number of tiles each containing a letter  $a$ ,  $b$ ,  $\dots$  or  $z$  from the alphabet. For any letter  $a$  from the alphabet your assignment includes  $N_a$  tiles the letter  $a$ . From the letters you are to construct any of the words  $w_1, \dots, w_n$ . This list might, for example, contain all words from a given dictionary. You may construct any word at most once and use any tile at most once. You receive  $v_j \geq 0$  points for making a word  $w_j$  and an additional bonus  $b_{ij} \geq 0$  for making both words  $w_i$  and  $w_j$ , for  $i, j \in \{1, 2, \dots, n\}$ .
- (a) Formulate a linear integer program to determine your optimal choice of words.
- (b) How does the formulation change if you are allowed to select 100 tiles with no restriction on your choice of the letters?
6. Consider the problem

$$\begin{aligned} \text{maximize } z &= x_1 + 2x_2 \\ \text{subject to} & \\ & x_1 + x_2 \leq 8 \\ & -x_1 + 2x_2 \leq 2 \\ & x_1 - x_2 \leq 4 \\ & x_2 \geq 0 \text{ and } x_2 \text{ integer} \\ & x_1 \in \{0, 1, 4, 6\} \end{aligned}$$

- (a) Reformulate the problem as an integer linear program.
- (b) How would your answer to part (a) change if the objective function was changed to maximize  $z = x_1^2 + 2x_2$ ?
7. Graph the following integer program

$$\begin{aligned} \text{maximize } z &= x_1 + 5x_2 \\ \text{subject to} & \\ & -4x_1 + 3x_2 \leq 6 \\ & 3x_1 + 2x_2 \leq 18 \\ & x_1, x_2 \geq 0 \text{ and } x_1, x_2 \text{ integer} \end{aligned}$$

Apply the branch-and-bound procedure, graphically solving each linear programming encountered. Interpret the branch-and-bound procedure graphically as well.