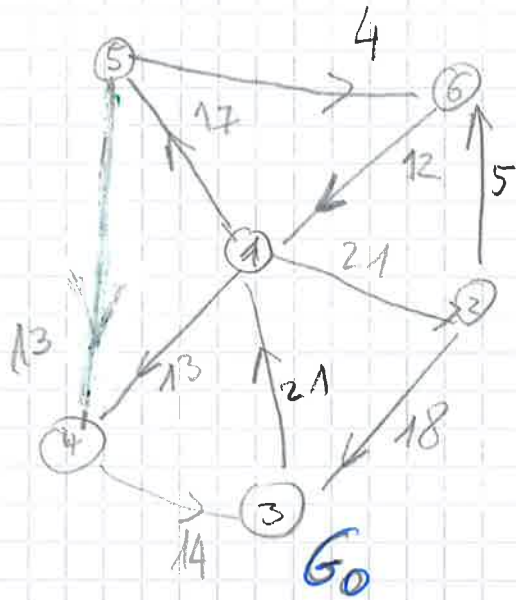


Example for the Edmonds branching algorithm ①



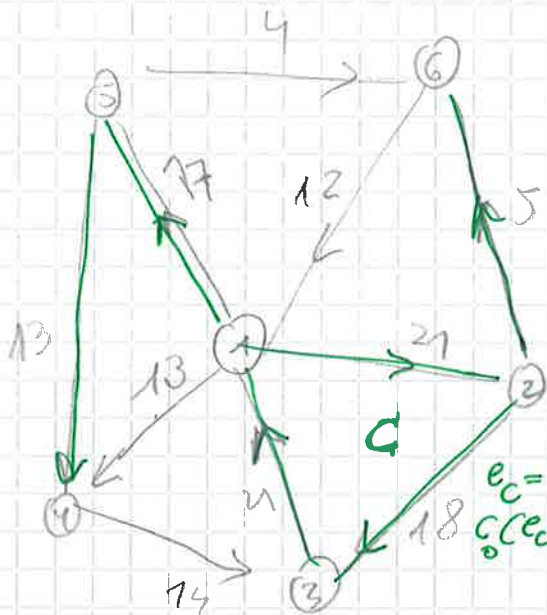
Step 1

$G = G_0 \quad i = 0 \quad C_0 = C$

Step 2: Construct B_0

Step 3 . B_0 contains a cycle $(1,2,3) = C$

$B := B_0$



B_0 Contract the circuits of G_0 and obtain G_1 with its weights C_1 and the representatives $\phi_1(\cdot)$

$e_C = 2,3$
 $C_0(e_C) = 18$

$\phi_1(5,6) = (5,6)$

$\phi_1(5,4) = (5,4)$

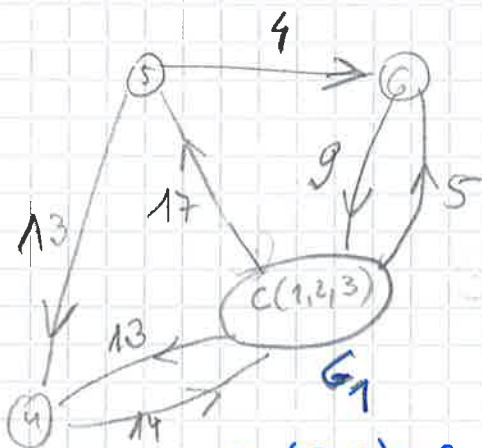
$\phi_1(6, C(1,2,3)) = (6,1)$

$\phi_1(C(1,2,3), 6) = (2,6)$

$\phi_1(C(1,2,3), 5) = (1,5)$

$\phi_1(C(1,2,3), 4) = (1,4)$

$\phi_1(4, C(1,2,3)) = (4,3)$



$C_1(5,6) = C_0(5,6) \quad C_1(5,4) = C_0(5,4)$

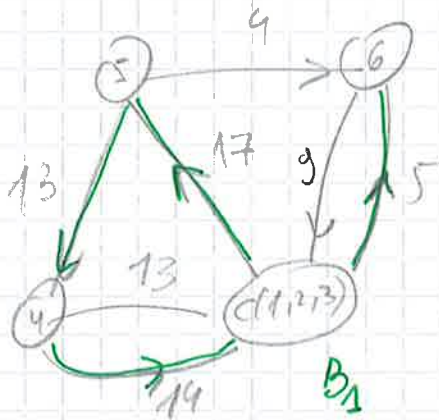
$C_1(C(1,2,3), 5) = C_0(1,5) \quad C_1(C(1,2,3), 6) = C_0(2,6)$

$C_1(C(1,2,3), 4) = C_0(1,4) = 13 \quad C_1(4, C(1,2,3)) = C_0(4,3) - C_0(2,3) + C_0(e_C) = 14$

$C_1(5, C(1,2,3)) = C_0(6,1) - C_0(3,1) + C_0(e_C) = 9$

Step 3: Second iterationSet $i=2$

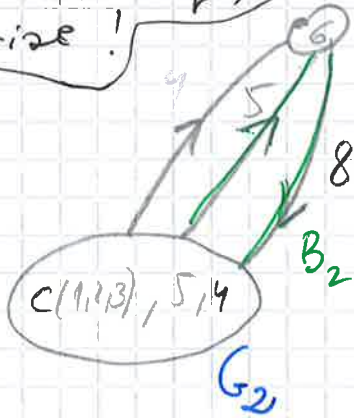
(2)

Construct B_1 

B_1 has a cycle
 $C((1,2,3), 5, 4)$

Contract it to obtain
 G_2 and its weights
 c_2 and the representatives
 $\phi_2(\cdot)$

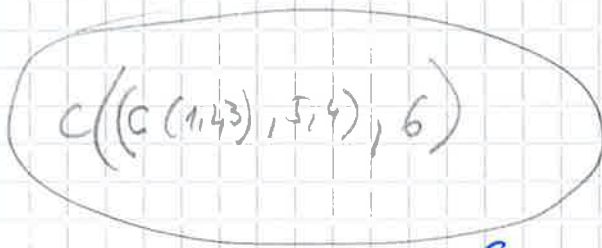
Notice
 parallel edges
 arise!

Step 3: third iterationset $i=2$ Construct B_2

B_2 has one cycle $C((C((1,2,3), 5, 4), 6))$

Contract it to obtain G_3

and the weights c_3 and
 the representatives
 $\phi_3(\cdot)$.

Step 3: fourth iterationset $i=3$ Construct B_3

Since $E(G_3) = \emptyset$ also $B_3 = \emptyset$.

So B_3 is cycle-free. Goto Step 5

set $B_1 = B_3$

$$i = 3 > 0$$

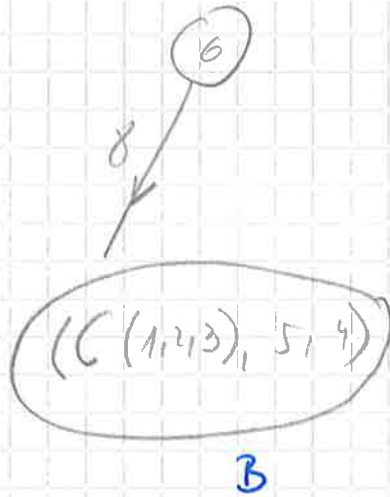
$$\text{let } B' = (V(G_2), \{\phi_3(e) : e \in \underbrace{E(B)}_{=\emptyset}\}) = (V(G_2), \emptyset)$$

$$\forall \text{ cycle } C \text{ in } B_2 \text{ do } [C = ((C(1,2,3), 5, 4), 6)]$$

$$E(B') = \underbrace{E(B')}_{=\emptyset} \cup (E(C) \setminus \{e_C\}) = (e_1, (C(1,2,3), 5, 4))$$

$$B = B'$$

$$i = 3 - 1 = 2$$



Step 5, second iteration

$$B' = (V(G_1), \{\phi_2(e) : e \in E(B)\})$$

$$\underbrace{\{C(1,2,3), 4, 5, 6\}}_{=\{e_1, C(1,2,3)\}}$$

\forall cycle C of B_1 (which is $(C(1,2,3), 5, 4)$)

if there is an edge entering the cycle

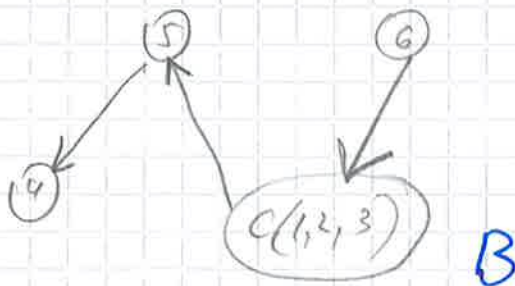
then

$$E(B') = E(B') \cup (E(C) \setminus \{e_1, C(1,2,3)\}) =$$

$$= \{(e_1, C(1,2,3)), (5, 4), (C(1,2,3), 5)\}$$

$$B = B'$$

$$i = 2 - 1 = 1$$



Step 5, third iteration

$$i = 1 > 0$$

$$B^1 = (V(B_0), \{ \phi_1(e) : e \in E(B) \})$$

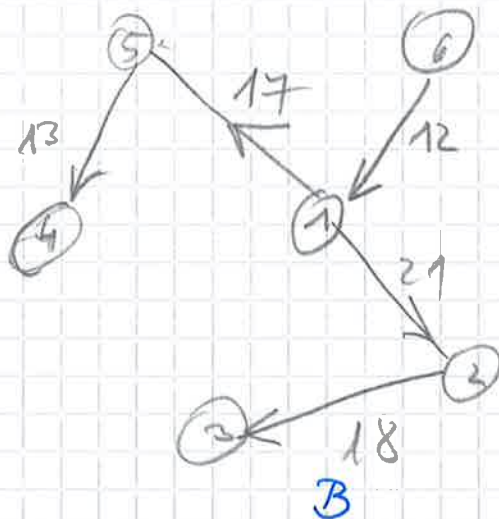
$$= (\{1,2,3,4,5,6\}, \{ (5,4), (1,5), (6,1) \})$$

if circuit C_i in B_0 (which is $C = C(1,2,3)$) do
 if there is an edge entering C (which is $(6,1)$)

$$\text{set } E(B^1) := E(B^1) \cup (E(C) \setminus \{(2,4)\})$$

$$= \{ (5,4), (1,5), (6,1), (1,2), (2,3) \}$$

$$\text{set } B = B^1 \quad i = i - 1 = 0$$



STOP

B is the maximum
 branching with
 weight $13 + 17 + 12 + 18 + 31$
 $= 81$