

Proof of Theorem 7.8

Assume w.l.o.g.  $k \geq 3$  ( $g(G) \geq 3$  holds for any graph).

Fix  $0 < \epsilon < \frac{1}{k}$  and set  $p := n^{\epsilon-1}$ .

Let  $X(G)$  denote the number of short cycles in a random graph  $G \in \mathcal{G}(n, p)$  (i.e. the number of cycles of length at most  $k$ ).

By Lemma 7.7 (on the exp.  $\mathbb{E}(X)$  of  $X$ ) we have

$$\begin{aligned} \underline{\mathbb{E}(X)} &= \sum_{i=3}^k \underbrace{\frac{\binom{n}{i}}{2i}} p^i \leq \frac{1}{2} \sum_{i=3}^k n^i p^i \stackrel{(*)}{\leq} \frac{1}{2} (k-2) n^k p^k \\ &= \mathbb{E}(X_i) \text{ where } X_i \text{ is the nr of } i\text{-cycles} \end{aligned}$$

The inequality  $(*)$  holds because  $n \cdot p = n n^{\epsilon-1} = n^\epsilon \geq 1$  and therefore  $(np)^i \leq (np)^k, \forall i \leq k$ .

Now by Lemma 7.6 (Markov's inequality) we have

$$\begin{aligned} \mathbb{P}\left(X \geq \frac{n}{2}\right) &\leq \frac{\mathbb{E}(X)}{\frac{n}{2}} \leq \frac{(k-2) n^k p^k}{\frac{n}{2}} = (k-2) n^{k-1} p^k \\ &= (k-2) n^{k-1} (n^\epsilon)^k = (k-2) n^{k\epsilon-1} \end{aligned}$$

By our choice of  $\epsilon \in (0, \frac{1}{k})$  we have  $\epsilon k - 1 < 0$

Thus  $\lim_{n \rightarrow \infty} \mathbb{P}\left(X \geq \frac{n}{2}\right) = 0$

Let now  $n$  be large enough such that  $\mathbb{P}\left(X \geq \frac{n}{2}\right) < \frac{1}{2}$

and  $\mathbb{P}\left(X \geq \frac{1}{2} \frac{n}{k}\right) < \frac{1}{2}$  (the latter is possible

by our choice of  $p$  in Lemma 7.7; here we had

$p \geq \frac{6k \ln n}{n}$  which is obviously fulfilled by  $p = \frac{n^\epsilon}{n} \neq \forall \epsilon > 0$ , if  $n$  is large enough)

Then there exists a graph  $G \in \mathcal{G}(n, p)$  with fewer than  $\frac{m}{2}$  short cycles and  $\chi(G) < \frac{1}{2} \frac{n}{k}$ .

From each of those cycles delete a vertex and let  $H$  be the graph resulting after this process.

Then  $|H| > \frac{n}{2}$  and  $H$  has no short cycles, so  $g(H) > k$ .

By the definition of  $G$  we get

$$\chi(H) \geq \frac{|H|}{\alpha(H)} \geq \frac{n/2}{\alpha(G)} > k \text{ and this completes the proof. } \square$$

Corollary: There are graphs with arbitrary large growth and arbitrary large values of the invariants  $k$ ,  $\epsilon$  and  $\delta$ .

Proof *Sketch, details to be worked out as a homework or exercise!*  
 We use a fact derived as a corollary of

$$\chi(G) \leq \text{col}(G) = \max \{ \delta(H) : H \subseteq G \} + 1 \text{ (Chapter 5)}$$

Corollary: Every graph  $G$  has a subgraph of minimum degree at least  $\chi(G) - 1$ .

This guarantees that  $\tilde{\mathcal{D}}(G)$  is large.

Further we use the Theorem of Roder

Th: Let  $0 \neq k \in \mathbb{N}$ . Every graph with  $d(G) \geq 4k$  has a  $(k+1)$ -connected subgraph  $H$  such that  $\epsilon(H) > \epsilon(G) - k$ .

Recall:  $\epsilon(G) := \frac{|\mathbb{E}(G)|}{|V(G)|}$ , the density of the graph

$$d(G) = \frac{1}{|V(G)|} \sum_{v \in V(G)} \text{deg}(v)$$

$k$  the connectivity number, i.e. the smallest number  $k$  s.t. removing  $k-1$  vertices does not disconnect  $G$  with  $V(G) \geq k$ .  $\square$

Properties of almost all graphs

Recall: a graph property is a class of graphs which is closed under isomorphism, i.e. one that contains with every graph  $G$  also the graphs isomorphic to  $G$ .

If  $p = p(n)$  is a fixed function (possibly constant) and  $\mathcal{P}$  is a graph property we may consider  $\lim_{n \rightarrow \infty} \mathbb{P}(G \in \mathcal{P})$  for  $G \in \mathcal{G}(n, p)$ .

If  $\lim_{n \rightarrow \infty} \mathbb{P}(G \in \mathcal{P}) = 1$  we say  $G \in \mathcal{P}$  for almost all  $G \in \mathcal{G}(n, p)$ , or that  $G \in \mathcal{P}$  almost surely in  $\mathcal{G}(n, p)$ .

If  $\lim_{n \rightarrow \infty} \mathbb{P}(G \in \mathcal{P}) = 0$  we say that almost no  $G \in \mathcal{G}(n, p)$

has property  $\mathcal{P}$  (example: we proved in Lemma 7.8 that for a certain  $p$ , almost no  $G \in \mathcal{G}(n, p)$  has a set of more than  $\frac{1}{2} \frac{n}{k}$  independent vertices)  $p \gg \frac{6k \ln n}{n}$

To illustrate this new concept we prove that for constant  $p$ , every fixed abstract graph  $H$  is an induced subgraph of almost all graphs.

Proposition 7.80 For every constant  $p \in (0, 1)$  and every graph  $H$ , almost every  $G \in \mathcal{G}(n, p)$  contains an induced copy of  $G$ .

Proof Let  $H$  be given and  $k := |H|$ .  $|H| = k$   
If  $n \geq k$  and  $U \subseteq \{0, \dots, n-1\}$  is a fixed set of  $k$  vertices of  $G$  then  $G[U]$  is isomorphic to  $H$

with a certain probability  $r > 0$ . This probability depends on  $p$  but not on  $n$  (why?)

Perfect graphs is  
 $r < 1$   $\Rightarrow$  nothing to show  
 $r > 1$   $\Rightarrow$  nothing to show

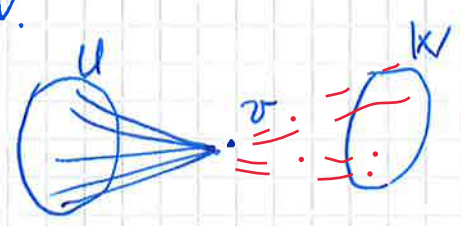
$G$  contains a collection of  $\lfloor \frac{n}{k} \rfloor$  disjoint such sets  $U$ . The probability that none of the  $G[U]$  is isomorphic to  $H$  is  $(1-r)^{\lfloor \frac{n}{k} \rfloor}$  because these events are independent, because the sets  $U$  are disjoint and hence also the edge sets in  $[U]^2$  are disjoint.

Thus  $P(H \not\subseteq G \text{ induced}) \leq (1-r)^{\lfloor \frac{n}{k} \rfloor} \xrightarrow{n \rightarrow \infty} 0$

( $r$  is fixed and also  $k$ )   
 prob that  $G[U] \cong H$  for any chosen set  $U$  among  $\lfloor \frac{n}{k} \rfloor$

The following lemma is a technical result needed to prove that many properties are almost sure in  $\mathcal{G}(n, p)$  for some suitably chosen  $p$ .

Given  $i, j \in \mathbb{N}$ , let  $P_{i,j}$  be the property that the considered graph contains, for any disjoint vertex sets  $U, W$  with  $|U| \leq i$  and  $|W| \leq j$  a vertex  $v \notin U \cup W$  that is adjacent to all vertices in  $U$  but none in  $W$ .



Lemma 7.16  $\forall p \in (0, 1)$  and  $i, j \in \mathbb{N}$  almost every graph  $G \in \mathcal{G}(n, p)$  has the property  $P_{i,j}$ .

Proof For fixed  $U, W$  and  $v \in G - (U \cup W)$ , the prob. that  $v$  is adjacent to all  $u \in U$  but to no  $w \in W$  is  $p^{|U|} q^{|W|} \geq p^i q^j$

$q = 1 - p$

Thus the probability that no such a suitable  $v$  event for these sets  $U$  and  $X$  is

$$(1-p^{|U|}q^{|W|})^{n-|U|-|W|} \leq (1-p^i q^j)^{n-i-j} \quad (\text{for } n \geq i+j),$$

since the corresponding events are independent for different  $v$ .  
 length  $(\dots \dots \dots |X|) \leq i+j$

There are at most  $n^{i+j}$  such sets  $U, W$  in  $V(G)$

Thus the probability that some such pair has no suitable  $v$  is at most

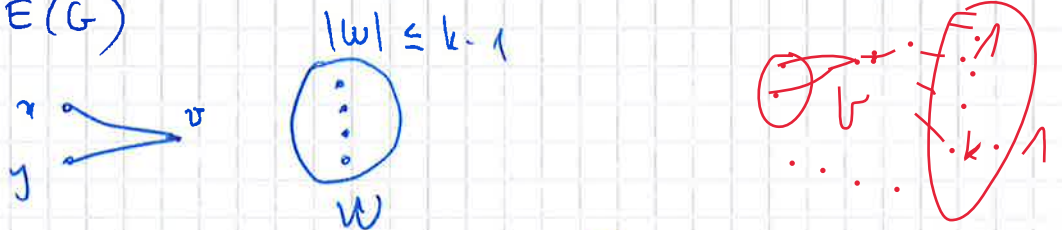
$$P(\cup_{U, W} E_{U, W}) \leq n^{i+j} (1-p^i q^j)^{n-i-j} \xrightarrow{n \rightarrow \infty} 0$$

$E_{U, W}$  event that no  $v$  with  $W$  exists

Corollary  $\forall p \in (0, 1), \forall k \in \mathbb{N}$ , almost every  $G \in \mathcal{G}(n, p)$  is  $k$ -connected.

Proof By lemma 7.10 above it is enough to show that every graph in  $\mathcal{P}_{2, k-1}$  is  $k$ -connected.

This is simple: Every graph  $G$  in  $\mathcal{P}_{2, k-1}$  is of order at least  $k+2$  and  $\nexists W \subseteq V(G)$  with  $|W| \leq k-1, \nexists x, y \notin W, \exists v \notin W \cup \{x, y\}$  with  $\{v, x\} \in E(G), \{v, y\} \in E(G)$



Thus  $W$  does not separate  $x$  from  $y$ .

Since this holds  $\nexists W \subseteq V(G)$  with  $|W| \leq k-1$  we conclude that  $G$  is  $k$ -connected.  $\square$

Remark: The hypothesis constructed in the proof of the Corollary has length 2, which is a stronger result than needed for the  $k$ -connectivity!