

## ( Chapter 3) Hamiltonicity: sufficient conditions

### Theorem 4.

(Chvátal, Erdős 1972)

Consider a graph  $G$  with  $n := |G| \geq 3$ . The inequality  $\kappa(G) \geq \alpha(G)$  implies the hamiltonicity of  $G$ .

( $\alpha(G)$  is the stability number of  $G$ , i.e.

$\alpha(G) := \max\{|S| : S \subseteq V(G) \text{ and } G[S] \text{ has no edges}\}.$ )

The result of Theorem 4 is best possible:

- (i) There exists a graph  $G$  with  $\kappa(G) = \alpha(G) - 1$  which is not Hamiltonian (e.g. the Petersen graph or  $K_{r,r+1}$ , as discussed in the practical)
- (ii) The statement also does not hold if  $\kappa(G)$  is replaced by  $\lambda(G)$ .

### Corollary 5.

There exists an  $O(mn)$  algorithm which either constructs a Hamiltonian cycle in  $G$  (with  $n := |V(G)$ ,  $m := |E(G)|$ ) or finds a stable set  $S$  and a separating set  $T$  in  $G$  such that  $|T| < |S|$  holds.

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### Theorem 6.

(Bondy, 1978)

Let  $G$  be a graph with  $|G| \geq 3$  and  $\deg(u) + \deg(v) \geq |G|$  for all  $u, v \in V(G)$  with  $\{u, v\} \notin E(G)$ . Then  $\kappa(G) \geq \alpha(G)$  holds.

## ( Chapter 3) Finding long cycles in graphs

The problem of determining the length of a longest cycle in a graph  $G$  is *NP*-hard.

We assume w.l.o.g. that  $G$  is 2-connected.

### Theorem 7.

*(Ore 1960, Bermond 1976, Linial 1976)*

*Let  $d \in \mathbb{N}$  and let  $G$  be a 2-connected graph such that  $\forall u, v \in V(G)$  with  $\{u, v\} \notin E(G)$ ,  $\deg(u) + \deg(v) \geq d$  holds. Then there exists a cycle of length at least  $\min\{n, d\}$  in  $G$ , where  $n := |G|$ .*

### Theorem 8.

*(Bauer, Broersma, Veldman and Rao 1989)*

*If  $G$  is a 2-connected graph with  $n$  vertices and connectivity  $\kappa(G)$  such that  $\deg(x) + \deg(y) + \deg(z) \geq n + \kappa(G)$  for any triple of independent vertices  $x, y, z \in V(G)$ , then  $G$  is Hamiltonian.*

Proof in *Journal of Combinatorial Theory, Series B* **47(2)**, 1989, 237–243.

## ( Chapter 3 ) Hamiltonian cycles and degree sequences

### Definition 1.

If  $G$  is a graph with  $n := |G|$  and vertex degrees  $d_1 \leq d_2 \leq \dots \leq d_n$ , then the  $n$ -tuple  $(d_1, d_2, \dots, d_n)$  is called **the degree sequence** of  $G$ . An arbitrary sequence of natural numbers  $(a_1, a_2, \dots, a_n)$  is called a **Hamiltonian sequence** if every graph with  $n$  vertices and degree sequence  $(d_1, d_2, \dots, d_n)$  which is pointwise not smaller than  $(a_1, a_2, \dots, a_n)$  (i.e.  $d_i \geq a_i, \forall i \in 1, n$ ), is Hamiltonian.

The degree sequence of a graph is unique while there may be different enumerations of the vertices all of them leading to the same degree sequence.

### Theorem 9.

(Chvátal 1972)

A sequence  $(a_1, a_2, \dots, a_n)$  of natural numbers with  $0 < a_1 \leq a_2 \leq \dots \leq a_n < n$  and  $n \geq 3$  is Hamiltonian iff  $a_i \leq i$  implies  $a_{n-i} \geq n - i$ , for every  $i < \frac{n}{2}$ .